Ruppert matrix as subresultant mapping

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09:29:35 (16)

-Ruppert matrix as subresultant mapping -

 $[\texttt{Prev}\Leftarrow] \; [\Rightarrow\texttt{Next}]$

Outline of this talk

- How to compute polynomial GCDs
- Sylvester matrix and Subresultant mapping
- Ruppert matrix and factoring polynomials
- Ruppert matrix for GCD of two polynomials
- Ruppert matrix for GCD of several polynomials
- Conclusion

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \Leftarrow] [\Rightarrow Next]$

How to compute polynomial GCDs

Euclidean Algorithm

Eculidean algorithm for polynomials f(x) and g(x) over reals. function GCD(f(x), g(x)) if g(x) = 0 return f(x)else return GCD(g(x), remainder of f(x) by g(x))

QR factoring of Sylvester matrix

The usual subresultant mapping is important. See the next slide...

Lots of other powerful approximate algorithms

-Ruppert matrix as subresultant mapping -	$[Prev \Leftarrow] [\Rightarrow Next]$
	-Ruppert matrix as subresultant mapping -

QR factoring of Sylvester matrix

Sylvester matrix of the following $f(x)$ and $g(x)$: $f(x) = f_n x^n + f_{n-1} x^{n-1} + \dots + f_1 x + f_0,$ $g(x) = g_m x^m + g_{m-1} x^{m-1} + \dots + g_1 x + g_0.$													
	(f_n)	f_{n-1}	•••	f_1	f_0	0	•••	0)					
	0	f_n	f_{n-1}	•••	f_1	f_0	·	:					
	•	·.	·.	·.	•••	·.	•.	0					
$\mathbf{S}_{\mathbf{r}}\mathbf{I}(\mathbf{f}_{\mathbf{r}}\mathbf{r})$	0		0	f_n	f_{n-1}	•••	f_1	f_0					
Syl(f, g) =	g_m	g_{m-1}	•••	g_1	g_0	0	•••	0					
	0	g_m	g_{m-1}	•••	g_1	g_0	•.	:					
	•	·.	·.	·	••••	·	·	0					
	0	•••	0	<i>g</i> _m	g_{m-1}	•••	g_1	g_0					
Well known relation between GCD and QR factoring of Sylvester													

The last non-zero row vector of *R* where Syl(f, g) = QR, is the coefficient vector of GCD of f(x) and g(x). (See: Laidacker, M.A., 1969 and others)

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[\texttt{Prev}\Leftarrow] \; [\Rightarrow\texttt{Next}]$

Subresultant mapping

Subresultant mapping S_r of $f_0(x)$ and $f_1(x)$ $S_r : \begin{cases} \mathcal{P}_{n_1 - r - 1} \times \mathcal{P}_{n_0 - r - 1} \longrightarrow \mathcal{P}_{n_0 + n_1 - r - 1} \\ (u_0, u_1) \longmapsto u_1 f_0 + u_0 f_1 \end{cases}$ where $\begin{cases} f_0(x) = f_{0,n_0} x^{n_0} + \dots + f_{0,1} x + f_{0,0} \\ f_1(x) = f_{1,n_1} x^{n_1} + \dots + f_{1,1} x + f_{1,0} \end{cases}$

• Sylvester subresultant matrix $S_r(f_0, f_1)$

 $S_r(f_0, f_1) = (C_{n_0-r}(f_1) \quad C_{n_1-r}(f_0))$

where $C_k(p)$ is a *k*-th convolution matrix of p(x)and $S_0(f_0, f_1)$ is the transpose of Sylvester matrix of $f_0(x)$ and $f_1(x)$.

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \leftarrow] [\Rightarrow Next]$

Convolution matrix

k-th convolution	n matri	x of po	lyno	mial p	(<i>x</i>).	
	(p_n)	0	•••	0	0)	
	p_{n-1}	p_n	·.	÷	÷	
	÷	p_{n-1}	•.	0	÷	
	p_0	÷	•.	p_n	0	
$C_k(p) =$	0	p_0	·.	p_{n-1}	p_n	
	÷	0	·.	÷	p_{n-1}	
	÷	÷	·	p_0	÷	
	0	0		0	p_0	
where	p(x)	$= p_n x'$	$n + \cdots$	$\cdot + p_1$	$x + p_0$.	

Subresultant mapping and Sylvester matrix

$$S_{r}: \begin{cases} \mathcal{P}_{n_{1}-r-1} \times \mathcal{P}_{n_{0}-r-1} \longrightarrow \mathcal{P}_{n_{0}+n_{1}-r-1} \\ (u_{0}, u_{1}) \longmapsto u_{1} f_{0} + u_{0} f_{1} \end{cases}$$
$$S_{r}(f_{0}, f_{1}) = (C_{n_{0}-r}(f_{1}) \quad C_{n_{1}-r}(f_{0})), \quad S_{0}(f_{0}, f_{1}) = {}^{t}Syl(f_{0}, f_{1})$$

|--|

Subresultant mapping and GCD

Well known fact (see Rupprecht, D. (1999) and so on)

(1) If r is the largest integer that

the subresultant mapping S_r is not injective,

we can compute the GCD of $f_0(x)$ and $f_1(x)$ from

the right null vector of $S_r(f_0, f_1)$. (2) The dimension of the null space is the degree of the polynomial GCD.

09:29:36 (16)

-Ruppert matrix as subresultant mapping

 $[Prev \Leftarrow] [\Rightarrow Next]$

Ruppert criterion

Reducibility of bivariate polynomials (Ruppert, W.M., 1999)



- Newton polytope ver. by Gao, S. and Rodrigues, V.M. (2003)
- Multivariate version by May, J.P. (2005)

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \Leftarrow] [\Rightarrow Next]$

Ruppert matrix R(f)

• The differential equation w.r.t. g(x, y) and h(x, y) \Rightarrow a linear equation w.r.t. unknown coefficients.

$$\begin{array}{rcl} \partial g & \partial f & \partial f & \partial h \\ f - - g - + h - - f - f - g = 0, \ g, \ h \in \mathbb{C}[x, y], \\ \partial y & \partial y & \partial x & \partial x \\ \deg_x g \leq & \deg_x f - 1, \ \deg_y g \leq & \deg_y f, \\ \deg_x h \leq & \deg_x f, \ \deg_y h \leq & \deg_y f - 2 \end{array}$$

$$R(f)^t (\text{ coefficient vectors of } g, h) = \emptyset$$

- SVD of Ruppert matrix: irreducibility radius.
- Ruppert matrix depends term orders: lexicographic in this talk.

Factoring with Ruppert and Gao's matrix

Known Algorithms using Corollary 1

Corollary 1

For a given $f(x, y) \in \mathbb{C}[x, y]$ that is square-free over $\mathbb{C}(y)$, the dimension (over \mathbb{C}) of the null space of R(f) is equal to "(the number of absolutely irreducible factors of f(x, y) over \mathbb{C}) -1".

and others by Kaltofen, Gao and so on.

• Correction: Corollary 1 from John May's thesis.

Kaltofen, E., May, J., Yang, Z., and Zhi, L.

Approximate factorization of multivariate polynomials using singular value decomposition. Manuscript, 22 pages. Submitted, Jan. 2006.

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \Leftarrow] [\Rightarrow Next]$

Purpose of this talk

GCD and Factorization

Primitive operations in Symbolic (and Numeric) computations.

Common Properties

Can be computed by some matrix decompositions.

Sylvester and Ruppert matices

Is there any relation between them?

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[\texttt{Prev}\Leftarrow] [\Rightarrow\texttt{Next}]$

How to link them

 $f(x, y) = f_0(x) + f_1(x) y$

 \blacksquare f(x, y) is reducible if $f_0(x)$ and $f_1(x)$ have a non-trivial GCD.

 \blacksquare f(x, y) is irreducible if they don't have any non-trivial GCD.

Coprimeness of $f_0(x)$ and $f_1(x)$ can be tested by a rank deficiency of their Sylvester matrix.

Irreducibility of $f(x, y) = f_0(x) + f_1(x) y$ can be tested by a rank deficiency of its Ruppert matrix.

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[Prev \Leftarrow] [\Rightarrow Next]$

Simple Result

Ι	Lemma 1
F	For any polynomials $f_0(x)$ and $f_1(x)$,
	the Sylvester matrix of $f_0(x)$ and $f_1(x)$
	and the Ruppert matrix of $f(x, y) = f_0(x) + f_1(x) y$
h	have the same information for computing their GCD,
	with the Ruppert's original differential equation and degree
C	constraints.
	The proof is in the proceedings

The proof is in the proceedings.

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \leftarrow] [\Rightarrow Next]$

Simple Result (Sylvester and Ruppert)

(fa!	5 fa4	f	a3	fa2	fa1	. fa	a 0	0	0	0		0)								
0	fa5	f	a4	fa3	fa2	2 fa	a 1	fa0	0	0	(0								
0	0	f	a5	fa4	fa3	5 fa	a 2	fa1	faC	0 ((0								
0	0		0	fa5	fa4	f fa	a 3	fa2	fa1	. fa0	(0								
0	0		0	0	fa	5 fa	a 4	fa3	fa2	2 fal	fa	a0								
fpi	5 fb4	f	b3	fb2	fb1	. fł	o0	0	0	0	(0								
0	fb5	£	b4	fb3	fb2	2 fl	5 1	fb0	0	0	(0								
0	0	f	b5	fb4	fb3	f f	5 2	fb1	fb0	0	(0								
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(0	0		0	0	fb5	5 fl	5 4	fb3	fb2	fb1	£	ьо)								
(0	fa5	0	fa	.4 0	fa	1 3	0	fa2	0	fa1	0	fa0	0	0	0	0	0	0	0	0
0	-fb5	0	-f]	b4 0	-f	b3	0	-fb2	0	-fb1	0	-fb0	0	0	0	0	0	0	0	0
0	0	0	fa	.5 0	fa	1 4	0	fa3	0	fa2	0	fa1	0	fa0	0	0	0	0	0	0
0	0	0	-f]	b5 0	-f	b4	0	-fb3	0	-fb2	0	-fb1	0	-fb0	0	0	0	0	0	0
0	0	0	0	0	fa	1 5	0	fa4	0	fa3	0	fa2	0	fal	0	fa0	0	0	0	0
0	0	0	0	0	-f	b5	0	-fb4	0	-fb3	0	-fb2	0	-fb1	0	-fb0	0	0	0	0
0	0	0	0	0) ()	0	fa5	0	fa4	0	fa3	0	fa2	0	fal	0	fa0	0	0
0	0	0	0	0) ()	0	-fb5	0	-fb4	0	-fb3	0	-fb2	0	-fb1	0	-fb0	0	0
0	0	0	0	0) ()	0	0	0	fa5	0	fa4	0	fa3	0	fa2	0	fal	0	fa0
0	0	0	0	0) ()	0	0	0	-fb5	0	-fb4	0	-fb3	0	-fb2	0	-fb1	0	-fb0

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[\texttt{Prev}\Leftarrow] \; [\Rightarrow\texttt{Next}]$

Alternative Result 1

Theorem 1

The polynomial GCD of $f_0(x)$ and $f_1(x)$ can be computed by Singular Value Decomposition (SVD) of Ruppert matrix of $f(x, y) = f_0(x) + f_1(x) y$ with the J.P.May's differential equation and degree constraints, if f(x, y) is square-free over $\mathbb{C}(y)$.

The proof is in the proceedings.

Difference between Ruppert and May's

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[Prev \Leftarrow] [\Rightarrow Next]$

Alternative Result 1(Ruppert with May)

1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	-fb4	fb5	0	0	0	0	0	0	0	0	0	0	0	0
	-2 fb3	0	2 £b5	0	0	0	0	0	0	0	0	0	0	0
	-3 fb2	-fb3	fb4	3 £b5	0	0	0	0	0	0	0	0	0	0
	-4 fb1	-2 fb2	0	2 fb4	4 fb5	0	0	0	0	0	0	0	0	0
	-5 fb0	-3fb1	-fb2	fb3	3 £b4	5 £b5	0	0	0	0	0	0	0	0
	0	-4 fb0	-2 fb1	0	2 fb3	4 fb4	0	0	0	0	0	0	0	0
	0	0	-3£b0	-fb1	fb2	3 £b3	0	0	0	0	0	0	0	0
	0	0	0	-2 fb0	0	2 fb2	0	0	0	0	0	0	0	0
	0	0	0	0	-fb0	fb1	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	-fa4	fa5	0	0	0	0	-fa5	0	0	0	fb5	0	0	0
	-2 fa3	0	2 fa5	0	0	0	-fa4	-fa5	0	0	fb4	fb5	0	0
	-3fa2	-fa3	fa4	3 fa5	0	0	-fa3	-fa4	-fa5	0	fb3	fb4	£b5	0
	-4 fa1	-2 fa2	0	2 fa4	4 fa5	0	-fa2	-fa3	-fa4	-fa5	fb2	fb3	fb4	fb5
	-5 fa0	-3fa1	-fa2	fa3	3 fa4	5 fa5	-fal	-fa2	-fa3	-fa4	fb1	fb2	£b3	fb4
	0	-4 fa0	-2 fa1	0	2 fa3	4 fa4	-fa0	-fal	-fa2	-fa3	fb0	fb1	fb2	fb3
	0	0	-3fa0	-fal	fa2	3 fa3	0	-fa0	-fal	-fa2	0	fb0	fb1	fb2
	0	0	0	-2 fa0	0	2 fa2	0	0	-fa0	-fal	0	0	fb0	fb1
	0	0	0	0	-fa0	fa1	0	0	0	-fa0	0	0	0	fb0

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[Prev \Leftarrow] [\Rightarrow Next]$

Alternative Result 2

Theorem 2

The polynomial GCD of $f_0(x)$ and $f_1(x)$ can be computed by QR factoring of the transpose of the last 3 n_0 rows of their Ruppert matrix $R(f) = R(f_0(x) + f_1(x) y)$ with J.P.May's equation.

The last non-zero row vector of the triangular matrix is the coefficient vector of their polynomial GCD.

The proof is in the proceedings.



Alternative Result 2(Last 3n₀ rows)

1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	-fb4	fb5	0	0	0	0	0	0	0	0	0	0	0	0
	-2 fb3	0	2 fb5	0	0	0	0	0	0	0	0	0	0	0
	-3 fb2	-fb3	fb4	3 fb5	0	0	0	0	0	0	0	0	0	0
	-4 fb1	-2 fb2	0	2 fb4	4 fb5	0	0	0	0	0	0	0	0	0
	-5fb0	-3fb1	-fb2	fb3	3 fb4	5 fb5	0	0	0	0	0	0	0	0
	0	-4 fb0	-2 fb1	0	2 fb3	4 fb4	0	0	0	0	0	0	0	0
	0	0	-3fb0	-fb1	fb2	3 £b3	0	0	0	0	0	0	0	0
	0	0	0	-2fb0	0	2 fb2	0	0	0	0	0	0	0	0
	0	0	0	0	-fb0	fb1	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	-fa4	fa5	0	0	0	0	-fa5	0	0	0	fb5	0	0	0
	-2fa3	0	2 fa5	0	0	0	-fa4	-fa5	0	0	fb4	fb5	0	0
	-3fa2	-fa3	fa4	3 fa5	0	0	-fa3	-fa4	-fa5	0	fb3	fb4	fb5	0
	-4 fal	-2fa2	0	2 fa4	4 fa5	0	-fa2	-fa3	-fa4	-fa5	fb2	fb3	fb4	fb5
	- 5 fa 0	-3fa1	-fa2	fa3	3 fa4	5 fa5	-fal	-fa2	-fa3	-fa4	fb1	fb2	fb3	fb4
	0	- 4 fa 0	-2 fa1	0	2 fa3	4 fa4	-fa0	-fal	-fa2	-fa3	fb0	fb1	fb2	fb3
	0	0	-3fa0	-fal	fa2	3 fa3	0	-fa0	-fal	-fa2	0	fb0	fb1	fb2
	0	0	0	- 2 fa 0	0	2 fa2	0	0	-fa0	-fal	0	0	fb0	fb1
	0	0	0	0	-fa0	fa1	0	0	0	-fa0	0	0	0	fb0

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[\texttt{Prev}\Leftarrow] \; [\Rightarrow\texttt{Next}]$

Generalized Sylvester matrix

09:2

Generalized Subresultant mapping S_r of $f_0(x),, f_k(x)$.								
$S_r: \begin{cases} \prod_{i=0}^k \mathcal{P}_{n_i-r-1} & \longrightarrow & \prod_{i=1}^k \mathcal{P}_{n_0+n_i-r-1} \\ \begin{pmatrix} u_0 \\ \vdots \\ u_k \end{pmatrix} & \longmapsto & \begin{pmatrix} u_1 f_0 + u_0 f_1 \\ \vdots \\ u_k f_0 + u_0 f_k \end{pmatrix}, n_i = \deg_x f_i(x) \end{cases}$								
Generalized Sylvester matrix $S_r(f_0,, f_k)$.								
$S_r(f_0, \ldots, f_k) =$	$(C_{n_0-r}(f_1))$ $C_{n_0-r}(f_2)$ \vdots $(C_{n_0-r}(f_k))$	$C_{n_1-r}(f_0)$ \emptyset	$ \begin{array}{c} 0\\ C_{n_2-r}(f_0)\\ \vdots\\ \dots\end{array} $	···· ·· Ø	$ \begin{bmatrix} \emptyset \\ \vdots \\ 0 \\ C_{n_k-r}(f_0) \end{bmatrix} $			
(see Rupprecht, D. (1999) and so on)								
If <i>r</i> is the largest integer that the subresultant mapping S_r is not injective, we can compute the GCD of $f_0(x),, f_k(x)$ from the right null vector of $S_r(f_0,, f_k)$.								
9:36 (16) -Ruppert matrix as subresultant mapping - [Prev⇔] [⇒Next]								

How to link them



Relation for several polynomials

Postponed as a future work:

J.P.May's differential equation becomes

$$\begin{cases}
f_0 g_i - g_0 f_i + \lambda_i (f_i f_0^{'} - f_0 f_i^{'}) = 0 \\
f_1 g_i - g_1 f_i + \lambda_i (f_i f_1^{'} - f_1 f_i^{'}) = 0 \\
\vdots \\
f_k g_i - g_k f_i + \lambda_i (f_i f_k^{'} - f_k f_i^{'}) = 0
\end{cases}$$
(*i* = 1, ..., *k*)

This means that

$$\begin{cases} u_0(x) = g_0(x) - \lambda_i f_0' \\ u_i(x) = -g_i(x) + \lambda_i f_i' \end{cases} \quad (i = 1, ..., k)$$

 \implies will be a part of proof of several polynomial version of theorem.

09:29:36 (16)	-Ruppert matrix as subresultant mapping -	$[Prev \Leftarrow] [\Rightarrow Next]$

Conclusion

Sylvester and Ruppert matrices

have some relations for computing GCDs.

Their relations lead to

no algorithm which computes GCDs efficiently.

The author hopes that

the relations in this talk will make some progress.

Thank you !!

09:29:36 (16)

-Ruppert matrix as subresultant mapping -

 $[Prev \Leftarrow] [\Rightarrow Next]$