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#### **Observations:**

- ▶ The solution T[1], T[2], T[3],... is completely determined by a set of boundary conditions that specify values for T[1],...,T[k].
- In fact, any k consecutive values completely determine the solution.
- k non-concecutive values might not be an appropriate set of boundary conditions (depends on the problem).

- First determine all solutions that satisfy recurrence relation.
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#### The solution space

$$S = \{ \mathcal{T} = T[1], T[2], T[3], \dots \mid \mathcal{T} \text{ fulfills recurrence relation } \}$$

is a vector space. This means that if  $T_1, T_2 \in S$ , then also  $\alpha T_1 + \beta T_2 \in S$ , for arbitrary constants  $\alpha, \beta$ .

How do we find a non-trivial solution?

We guess that the solution is of the form  $\lambda^n$ ,  $\lambda \neq 0$ , and see what happens. In order for this guess to fulfill the recurrence we need

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This means that if  $\lambda_i$  is a root (Nullstelle) of  $P[\lambda]$  then  $T[n]=\lambda_i^n$  is a solution to the recurrence relation.

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#### Lemma 1

Assume that the characteristic polynomial has k distinct roots  $\lambda_1, \ldots, \lambda_k$ . Then all solutions to the recurrence relation are of the form

$$\alpha_1\lambda_1^n + \alpha_2\lambda_2^n + \cdots + \alpha_k\lambda_k^n$$
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#### Proof

There is one solution for every possible choice of boundary conditions for  $T[1], \ldots, T[k]$ .

We show that the above set of solutions contains one solution for every choice of boundary conditions.



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$$\alpha_{1} \cdot \lambda_{1} + \alpha_{2} \cdot \lambda_{2} + \cdots + \alpha_{k} \cdot \lambda_{k} = T[1]$$

$$\alpha_{1} \cdot \lambda_{1}^{2} + \alpha_{2} \cdot \lambda_{2}^{2} + \cdots + \alpha_{k} \cdot \lambda_{k}^{2} = T[2]$$

$$\vdots$$

$$\alpha_{1} \cdot \lambda_{1}^{k} + \alpha_{2} \cdot \lambda_{2}^{k} + \cdots + \alpha_{k} \cdot \lambda_{k}^{k} = T[k]$$



#### Proof (cont.).

$$\begin{pmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_k^2 \\ & & \vdots & & \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_k^k \end{pmatrix} \begin{pmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_k \end{pmatrix} = \begin{pmatrix} T[1] \\ T[2] \\ \vdots \\ T[k] \end{pmatrix}$$



#### Proof (cont.).

Suppose I am given boundary conditions T[i] and I want to see whether I can choose the  $\alpha_i's$  such that these conditions are met:

$$\begin{pmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_k^2 \\ & \vdots & & \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_k^k \end{pmatrix} \begin{pmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_k \end{pmatrix} = \begin{pmatrix} T[1] \\ T[2] \\ \vdots \\ T[k] \end{pmatrix}$$

We show that the column vectors are linearly independent. Then the above equation has a solution.



$$\begin{vmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_{k-1} & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_{k-1}^2 & \lambda_k^2 \\ \vdots & \vdots & & \vdots & \vdots \\ \lambda^k & \lambda^k & \cdots & \lambda^k & \lambda^k \end{vmatrix}$$

$$\begin{vmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_{k-1} & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_{k-1}^2 & \lambda_k^2 \\ \vdots & \vdots & & \vdots & \vdots \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_{k-1}^k & \lambda_k^k \end{vmatrix} = \prod_{i=1}^k \lambda_i \cdot \begin{vmatrix} 1 & 1 & \cdots & 1 & 1 \\ \lambda_1 & \lambda_2 & \cdots & \lambda_{k-1} & \lambda_k \\ \vdots & \vdots & & \vdots & \vdots \\ \lambda_1^{k-1} & \lambda_2^{k-1} & \cdots & \lambda_{k-1}^{k-1} & \lambda_k^{k-1} \end{vmatrix}$$

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$$=\prod_{i=1}^k \lambda_i \cdot \begin{vmatrix} 1 & \lambda_1 & \cdots & \lambda_1^{k-2} & \lambda_1^{k-1} \\ 1 & \lambda_2 & \cdots & \lambda_2^{k-2} & \lambda_2^{k-1} \\ \vdots & \vdots & & \vdots & \vdots \\ 1 & \lambda_k & \cdots & \lambda_k^{k-2} & \lambda_k^{k-1} \end{vmatrix}$$

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$$\begin{vmatrix} 1 & \lambda_{1} - \lambda_{1} \cdot 1 & \cdots & \lambda_{1}^{k-2} - \lambda_{1} \cdot \lambda_{1}^{k-3} & \lambda_{1}^{k-1} - \lambda_{1} \cdot \lambda_{1}^{k-2} \\ 1 & \lambda_{2} - \lambda_{1} \cdot 1 & \cdots & \lambda_{2}^{k-2} - \lambda_{1} \cdot \lambda_{2}^{k-3} & \lambda_{2}^{k-1} - \lambda_{1} \cdot \lambda_{2}^{k-2} \\ \vdots & \vdots & & \vdots & & \vdots \\ 1 & \lambda_{k} - \lambda_{1} \cdot 1 & \cdots & \lambda_{k}^{k-2} - \lambda_{1} \cdot \lambda_{k}^{k-3} & \lambda_{k}^{k-1} - \lambda_{1} \cdot \lambda_{k}^{k-2} \end{vmatrix}$$

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$$\begin{vmatrix} 1 & 0 & \cdots & 0 & 0 \\ 1 & (\lambda_2 - \lambda_1) \cdot 1 & \cdots & (\lambda_2 - \lambda_1) \cdot \lambda_2^{k-3} & (\lambda_2 - \lambda_1) \cdot \lambda_2^{k-2} \\ \vdots & \vdots & & \vdots & \vdots \\ 1 & (\lambda_k - \lambda_1) \cdot 1 & \cdots & (\lambda_k - \lambda_1) \cdot \lambda_k^{k-3} & (\lambda_k - \lambda_1) \cdot \lambda_k^{k-2} \end{vmatrix}$$



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$$\begin{vmatrix} \sum_{i=2}^{k} (\lambda_i - \lambda_1) \cdot \begin{pmatrix} 1 & \lambda_2 & \cdots & \lambda_2^{k-3} & \lambda_2^{k-2} \\ \vdots & \vdots & & \vdots & \vdots \\ 1 & \lambda_k & \cdots & \lambda_k^{k-3} & \lambda_k^{k-2} \end{pmatrix}$$



Repeating the above steps gives:

$$\begin{vmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_{k-1} & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_{k-1}^2 & \lambda_k^2 \\ \vdots & \vdots & & \vdots & \vdots \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_{k-1}^k & \lambda_k^k \end{vmatrix} = \prod_{i=1}^k \lambda_i \cdot \prod_{i>\ell} (\lambda_i - \lambda_\ell)$$

Hence, if all  $\lambda_i$ 's are different, then the determinant is non-zero.



#### What happens if the roots are not all distinct?

Suppose we have a root  $\lambda_i$  with multiplicity (Vielfachheit) at least 2. Then not only is  $\lambda_i^n$  a solution to the recurrence but also  $n\lambda_i^n$ .

To see this consider the polynomia

$$P[\lambda] \cdot \lambda^{n-k} = c_0 \lambda^n + c_1 \lambda^{n-1} + c_2 \lambda^{n-2} + \dots + c_k \lambda^{n-k}$$

Since  $\lambda_i$  is a root we can write this as  $Q[\lambda] \cdot (\lambda - \lambda_i)^2$ . Calculating the derivative gives a polynomial that still has root  $\lambda_i$ .



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Since  $\lambda_i$  is a root we can write this as  $Q[\lambda] \cdot (\lambda - \lambda_i)^2$ . Calculating the derivative gives a polynomial that still has root  $\lambda_i$ .



#### This means

$$c_0n\lambda_i^{n-1}+c_1(n-1)\lambda_i^{n-2}+\cdots+c_k(n-k)\lambda_i^{n-k-1}=0$$

Hence,

$$c_0 \underbrace{n\lambda_i^n}_{T[n]} + c_1 \underbrace{(n-1)\lambda_i^{n-1}}_{T[n-1]} + \cdots + c_k \underbrace{(n-k)\lambda_i^{n-k}}_{T[n-k]} = 0$$



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Suppose  $\lambda_i$  has multiplicity j. We know that

$$c_0 n \lambda_i^n + c_1 (n-1) \lambda_i^{n-1} + \dots + c_k (n-k) \lambda_i^{n-k} = 0$$

(after taking the derivative; multiplying with  $\lambda$ ; plugging in  $\lambda_i$ )

Doing this again gives

$$c_0 n^2 \lambda_i^n + c_1 (n-1)^2 \lambda_i^{n-1} + \dots + c_k (n-k)^2 \lambda_i^{n-k} = 0$$

We can continue j-1 times.

Hence,  $n^\ell \lambda_i^n$  is a solution for  $\ell \in {0,\ldots,j-1}.$ 



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#### Lemma 2

Let  $P[\lambda]$  denote the characteristic polynomial to the recurrence

$$c_0T[n] + c_1T[n-1] + \cdots + c_kT[n-k] = 0$$

Let  $\lambda_i$ ,  $i=1,\ldots,m$  be the (complex) roots of  $P[\lambda]$  with multiplicities  $\ell_i$ . Then the general solution to the recurrence is given by

$$T[n] = \sum_{i=1}^{m} \sum_{j=0}^{\ell_i - 1} \alpha_{ij} \cdot (n^j \lambda_i^n) .$$

The full proof is omitted. We have only shown that any choice of  $\alpha_{ij}$ 's is a solution to the recurrence.



$$T[0] = 0$$
  
 $T[1] = 1$   
 $T[n] = T[n-1] + T[n-2]$  for  $n \ge 2$ 

The characteristic polynomial is

$$\lambda^2 - \lambda - 1$$

Finding the roots, gives

$$\lambda_{1/2} = \frac{1}{2} \pm \sqrt{\frac{1}{4} + 1} = \frac{1}{2} \left( 1 \pm \sqrt{5} \right)$$



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$$\alpha\left(\frac{1+\sqrt{5}}{2}\right)+\beta\left(\frac{1-\sqrt{5}}{2}\right)=1 \Rightarrow \alpha-\beta=\frac{2}{\sqrt{5}}$$

Hence, the solution is

$$\frac{1}{\sqrt{5}} \left[ \left( \frac{1+\sqrt{5}}{2} \right)^n - \left( \frac{1-\sqrt{5}}{2} \right)^n \right]$$



Consider the recurrence relation:

$$c_0T(n) + c_1T(n-1) + c_2T(n-2) + \cdots + c_kT(n-k) = f(n)$$

with  $f(n) \neq 0$ .

While we have a fairly general technique for solving homogeneous, linear recurrence relations the inhomogeneous case is different.



The general solution of the recurrence relation is

$$T(n) = T_h(n) + T_p(n) ,$$

where  $T_h$  is any solution to the homogeneous equation, and  $T_p$  is one particular solution to the inhomogeneous equation.

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#### Example:

$$T[n] = T[n-1] + 1$$
  $T[0] = 1$ 

Then,

$$T[n-1] = T[n-2] + 1 \qquad (n \ge 2)$$

Subtracting the first from the second equation gives,

$$T[n] - T[n-1] = T[n-1] - T[n-2]$$
  $(n \ge 2)$ 

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$$T[n] = 2T[n-1] - T[n-2]$$
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 gives  $1 + \beta = 2 \Longrightarrow \beta = 1$ .

If f(n) is a polynomial of degree r this method can be applied r+1 times to obtain a homogeneous equation:

$$T[n] = T[n-1] + n^2$$

Shift

$$T[n-1] = T[n-2] + (n-1)^2 - T[n-2] + n^2 - 2n + 1$$

$$T[n] - T[n-1] = T[n-1] - T[n-2] + 2n - 1$$

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Difference:

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and so on...

