Technische Universität München Fakultät für Informatik Lehrstuhl für Effiziente Algorithmen Prof. Dr. Ernst W. Mayr Chris Pinkau

# **Complexity Theory**

## Due date: May 26, 2014 before class!

#### Problem 1 (10 Points)

Show that  $\mathbf{SPACE}(n) \neq \mathcal{NP}$ . (Note that it is unknown if either class is contained in the other.)

#### Problem 2 (10 Points)

Define the class  $\mathbf{E} = \bigcup_{c} \mathbf{DTIME}(2^{cn}).$ 

- 1. Is **E** closed under polynomial-time reductions?
- 2. Show that  $\mathcal{P}^{\mathbf{E}} = \mathbf{E}\mathbf{X}\mathbf{P}$ .

### Problem 3 (10 Points)

- 1. Argue that at least one of the assumptions  $\mathbf{L} \neq \mathcal{P}$  and  $\mathcal{P} \neq \mathbf{PSPACE}$  is true.
- 2. Use padding to show that if  $\mathcal{P} = \mathbf{L}$ , then  $\mathbf{EXP} = \mathbf{PSPACE}$ .

#### Problem 4 (10 Points)

Consider the problem of checking a boolean formula's syntactical correctness. Show that this problem can be decided in log-space, even if we have no precedence relation between the boolean operators and force precedence behavior with parentheses, e.g.  $(x \wedge y) \vee (\overline{z} \wedge x) \vee \overline{y} \vee z$  is a valid formula, as is  $(x \wedge (y \vee \overline{z}) \wedge x) \vee \overline{y} \vee z$ , while  $x \wedge y \vee \overline{z} \wedge x \vee \overline{y} \vee z$ is not.